

#### VSAM Record Level Sharing (RLS Overview) Part 1 and 2

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# Agenda

- Overview of RLS z/OS Release enhancements.
- IBM Products Exploiting RLS.

Record Level Sharing - Design direction.

- Review of base VSAM.
  - •Share Options
  - Buffering
  - locking
  - •RAS
  - •Performance Measurements
- Review of RLS
  - Share Options
  - Buffering
  - locking
  - •RAS
  - Performance Measurement

# Agenda (continued)

RLS/TVS Configuration Changes

Parmlib Changes

SYSPLEX with SMSVSAM

SMSVSAM Initialization

SMSVSAM Commands

RLS/CICS Environment

- •CICS and base VSAM FOR configuration
- •CICS and RLS configuration
- •RLS/CICS data recovery
- •RLS/CICS automation enhancements

# Agenda (continued)

Transactional VSAM (TVS)

•Hardware/Software Requirements

•Application Requirements

Multiple Lock Structure (future enhancement)

Recommended APARs

Summary

RLS z/OS Release Enhancements

### RLS z/OS Release Enhancements

- OS/390 2.1 VSAM RLS general availability (1996)
- z/OS 1.4 Transactional VSAM (priced feature)
- All z/OS Releases RAS support shipped via APARs
- z/OS 1.7 VSAM RLS 64 Buffering
- z/OS 1.8 RMF support for 64 bit buffering. RAS support.
   RSM changes.
- z/OS 1.9 RAS support, sysplex wide dumping.
- z/OS x.x Multiple Lock Structure support
- z/OS x.x CA Reclaim

# IBM Products Exploiting VSAM RLS

# IBM Products Exploiting RLS/TVS:

CICS

∎HSM

INFOMAN

SCLM

■IMS (RLS and TVS)

Record Level Sharing (RLS) – Design Direction

# Record Level Sharing (RLS) - Design

VSAM RLS is another method of access, to your existing VSAM files, which provides full read and write integrity at the record level, to any number of users in your parallel sysplex.

# **Review of Base VSAM**

### **Review of Base VSAM**

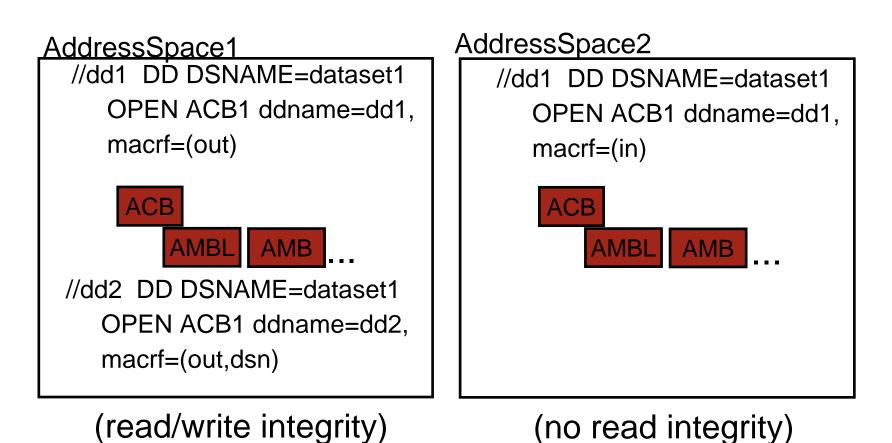
- Share options
- Buffering
- Locking
- RAS
- Performance Measurements

### **Review of Base VSAM**

#### Share options.

- attribute of the data set.
- SHAREOPTIONS(crossregion,crosssystem)
  - •SHAREOPTIONS(1,x) Defined as one user opened to the data set for read/write or any number of users for input only. VSAM provides full read/write integrity.
  - •SHAREOPTIONS(2,x) Defined as one user opened to the data set for read/write and any number of users for input VSAM provides full read/write integrity for the read/write user, however, the readers do not receive read integrity.
  - •SHAREOPTIONS(3,x) Defined as any number of users opened to the data set for read/write. VSAM does not provide any read/write integrity.
  - •SHAREOPTIONS(4,x) VSAM will flush buffers after each request.
- ACB MACRF=(DDN/DSN) is the only real mechanism for sharing VSAM files.

# Example of ShareOptions (2,x)



### **Base VSAM - Buffering**

 Base VSAM provides 3 types of buffering: ACB macrf=(NSR/LSR/GSR).

•NSR - Non-Shared Resources

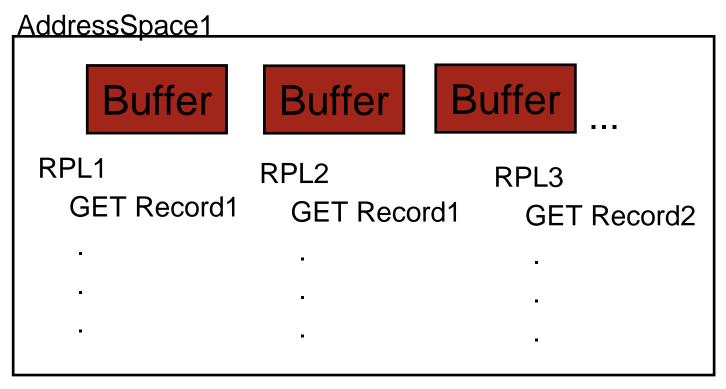
•LSR - Local Shared Resources

•GSR - Global Shared Resources

•For LSR/GSR, user defined the buffer pool:

POOL1 BLDVRP BUFFERS=(1024(5)), STRNO=4, TYPE=LSR, MODE=31, RMODE31=ALL

# Example of LSR Buffering



(read/write integrity)

### **Base VSAM - Locking**

- Base VSAM serializes on a CI level.
- Multiple users attempting to access the same CI for read and write either defer on the CI or are returned an exclusive control conflict error by VSAM.
- CIs with many records per CI, or applications that repeatedly access the same CI can have a performance impact due to retrying of exclusive control conflict errors.

#### Example of Base VSAM LSR Serialization

Scope = Single LSR Buffer Pool Granularity = Control Interval Ownership = RPL

GET UPD RPL\_1

(Record B)

GET UPD RPL\_2

(Record E)

fails - Exclusive Control Conflict



Control Interval

### Base VSAM - RAS

Base VSAM has little to no first time data capture, and internal recovery, for logic errors.

- All resources are obtained in a single address space.
- EOT acted as cleanup routine (plus estae stacked by open/close).
- Performance highly valued over RAS.
- RAS in general was not a major requirement when VSAM was developed.

End result:

- Difficult problems to debug.
- •Broken data sets and data integrity problems.

### Base VSAM – Performance Measurements

Base VSAM provides SMF 62 and 64 records.

- SMF 62 Created by OPEN for each ACB.
- SMF 64 Created by EOV and CLOSE for each ACB, however, the stats represent the sum of all ACBs connected to the control block structure.

# **Review of RLS**

# Review of RLS

#### Share options

Example of RLS Readers/Writers

Example of Shareoption (2,x) with RLS and base VSAM

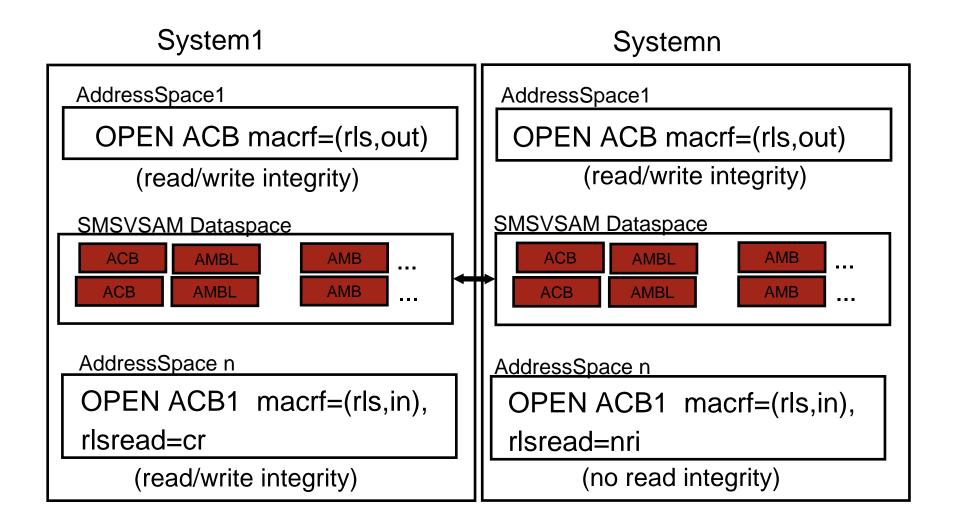
- Buffering
- Locking
- RAS
- Performance Measurements

## Review of RLS

#### Share options.

- largely ignored by RLS.
- Exception is SHAREOPTIONS(2,x) -
  - •Now defined as one user opened to the data set for non-RLS read/write and any number of users for non-RLS read. VSAM provides full read/write integrity for the non-RLS read/write user, however, the readers do not receive read integrity.
  - Or, any number of users opened for RLS read/write and any number of users for non-RLS read. VSAM provides full read/write integrity for the RLS users and no read integrity for the non-RLS readers.

# Example of RLS Readers/Writers



# Example of Shareoption (2,x) with RLS and base VSAM

System1 Systemn AddressSpace1 AddressSpace1 OPEN ACB macrf=(rls,out) OPEN ACB macrf=(rls,out) (read/write integrity) (read/write integrity) **SMSVSAM** Dataspace **SMSVSAM** Dataspace ACB AMB ACB AMBL AMB AMBL ACB AMB ACB AMBL AMBL AMB AddressSpace2 AddressSpace2 OPEN ACB1 macrf=(rls,in), OPEN ACB1 macrf=(nsr,in) ACB AMBL AMB rlsread=cr (no read integrity) (read/write integrity)

# **RLS** - Buffering

VSAM now provides 4 types of buffering: ACB macrf=(NSR/LSR/GSR/RLS).

- •NSR Non-Shared Resources
- •LSR Local Shared Resources
- •GSR Global Shared Resources
- •RLS Record Level Sharing
- Each image in the sysplex has one 31 bit local buffer pool, (located in a dataspace) with a current maximum size of 1.7 gig and one 64 bit pool located in the SMSVSAM address space. Both buffer pools are managed by LRU.
- Pool sizes controlled by PARMLIB parameters: RLS\_Max\_Pool\_Size (31 bit pool) and RLS\_Max\_Pool\_Size (31 bit pool) and
- RLSAboveTheBarMaxPoolSize (64 bit pool).
- Buffer coherency is maintained through the use of CF cache structures and the XCF cross-invalidation function.

The LRU for the 31 bit pool operates in the following 4 modes:

- Normal Mode Total pool size is less than 80% of RLS\_Max\_Pool\_Size.
- Maintenance Mode Total pool size is greater than 80% and less than 120% of RLS\_Max\_Pool\_Size.
- Accelerated Mode Total pool size is greater than 120% and less than 2\* RLS\_Max\_Pool\_Size.
- Panic Mode Total pool size is greater than 2\* RLS\_Max\_Pool\_Size or greater than 1728M.

#### The LRU will release 31 bit buffers as follows:

- Normal Mode IGWBLCRU will release invalid and paged out buffers.
  Initial Free\_UIC = 240.
  - •Buffer\_UIC + 1.

•Maximum age of buffers is 60 minutes.

- Maintenance Mode Reduce Initial\_Free\_UIC by 1. If Buffer\_UIC > Intial\_Free\_UIC\_Count then buffer is released (22.5 minutes max).
- •Accelerated Mode Reduce Initial\_Free\_UIC by 4. If Buffer\_UIC > Initial\_Free\_UIC then buffer is released. Requests for new buffers will first be stolen. If there are no buffers to steal a new get block will be done (7.5 minutes max).
- •Panic Mode Reduce Initial\_Free\_UIC by 8. If Buffer\_UIC > Initial\_Free\_UIC then buffer is released. Requests for new buffers will first be stolen (3.75 minutes max). If no buffers to steal, the request will be put to sleep until the LRU runs.

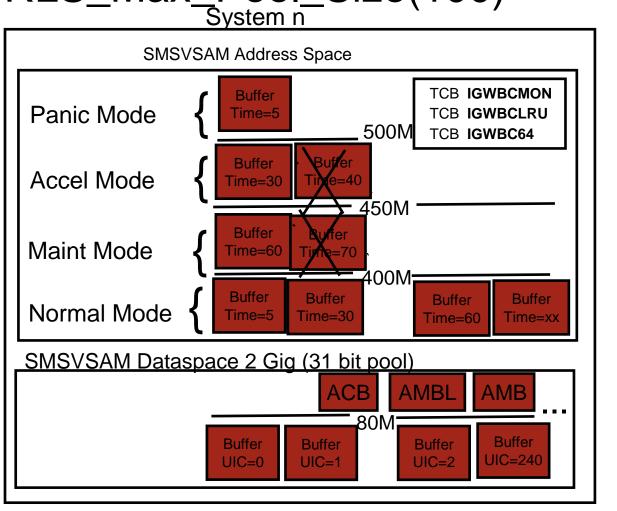
#### Setting the Local Buffer Pool Size – Considerations (cont):

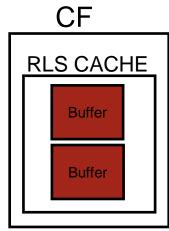
- •The LRU for the 64 bit buffer pool operates in four modes:
  - •Normal Mode Total 64 bit pool size is less than 80% of RLSAboveTheBarMaxPoolSize.
  - •Maintenance Mode Total 64 bit pool size is greater than 80% and less than 90% of RLSAboveTheBarMaxPoolSize.
  - •Accelerated Mode Total 64 bit pool size is greater than 90% and less than 100% of RLSAboveTheBarMaxPoolSize.
  - •Panic Mode Total 64 bit pool size is greater than 100% of RLSAboveTheBarMaxPoolSize

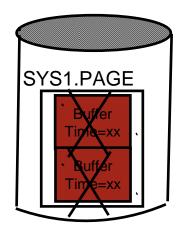
#### The LRU will release 64 bit buffers as follows:

- •Normal Mode Buffers 60 minutes or older will be released.
- Maintenance Mode Buffers 60 minutes or older will be released.
- Accelerated Mode Buffers 30 minutes are older will be released.
   Requests for new buffers will first be stolen. If there are no buffers to steal a new get block will be done.
- Panic Mode Buffers 5 minutes are older will be released. Requests for new buffers will first be stolen. If there are no buffers to steal, the request will sleep until LRU runs.

# RLSAboveTheBarMaxPoolSize(500) RLS\_Max\_Pool\_Size(100)







# Setting up Parameters/Structures sizes

Local Buffer Pool Sizes:

- •RLS\_MAX\_POOL\_SIZE(nnnn) Where nnnn = (10 to 9999), anything over 1500 is treated as a maximum of 1728M.
  •RLSAboveTheBarMaxPoolSize(sysname1,nnnn) Where nnnn is either 0, or 500M to
  - •RLSAboveTheBarMaxPoolSize(sysname1,nnnn) Where nnnn is either 0, or 500M to 2,000,000M
  - •RLS\_MaxCFFeatureLevel(Z/A)
- Pool Size values are a goal for which the LRU tries to maintain. If more buffers are required at any given time, the pool may temporarily exceed the values set.
- Real Storage Total amount of buffer pools should not exceed amount of real storage. A paged out buffer is immediately freed by the LRU.

## Sizing the RLS Cache Structures

The "ideal" cache structure size:

- Total\_Cache\_Sturcture\_sizes = ((RLS\_Max\_Pool\_Size) \* Number\_of\_SMSVSAMs\_in\_Sysplex) + (RLSAboveTheBarMaxPoolSize(system1) + ... +RLSAboveTheBarMaxPoolSize(systemn))
- Assumes the following:
  - RLS\_MaxCFFeaturelevel(A) caching all data
  - No sharing of data across the sysplex.
  - If more than one cache structure to be allocated, Data sets are "evenly" distributed (size, number, amount of data accessed) between the individual cache structures.

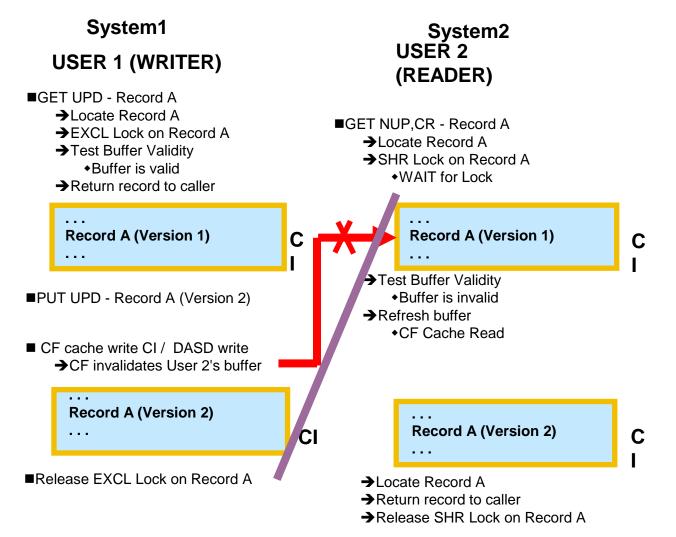
# Example CPU Time for GET Request

- Get request in which all CIs were found in the local buffer
   pool: .0001xx .0002xx seconds
- Get Request in which at least the one CI is read from DASD: .001x - .02xxxx Seconds

## Example CPU Time for GET Request

- Get request in which all CIs were found in the local buffer
   pool: .0001xx .0002xx seconds
- Get Request in which at least the one CI is read from DASD: .001x - .02xxxx Seconds

#### **RLS Buffer Invalidate Example**



## **RLS** - Locking

RLS serializes on a record level.

- Users updating or inserting a record will hold the lock exclusive for the duration of the write request or transaction.
- Users reading a record will hold the lock share when consistent read (CR) is specified. Lock is released at end of request

•ACB RLSREAD=CR

•//dd1 DD dsn=datasetname,RLS=CR

## RLS - Locking (cont.)

Users reading a record will not obtain any locks when no read intergrity (NRI) is specified.

•ACB RLSREAD=NRI

•//dd1 DD dsn=datasetname,RLS=NRI

Users reading a record will hold the lock share when consistent read extended (CRE) is specified. The lock is released at the end of the transaction:

•ACB RLSREAD=CRE

//dd1 DD dsn=datasetname,RLS=CRE

RLS locking is performed through the use of a CF lock structure and the XES locking services.

#### Example of VSAM RLS Serialization

Scope = Sysplex Granularity = Record Ownership = CICS Transaction or Batch Job

CICS1.Tran1

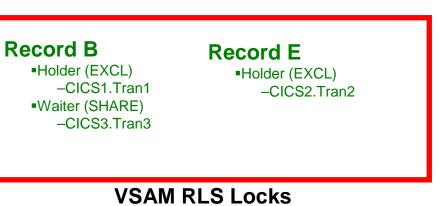
CICS2.Tran2

GET UPD RPL\_1 ( Record B) GET UPD RPL\_2 (Record E)

GET CR RPL\_3 (Record B) –Waits for record lock

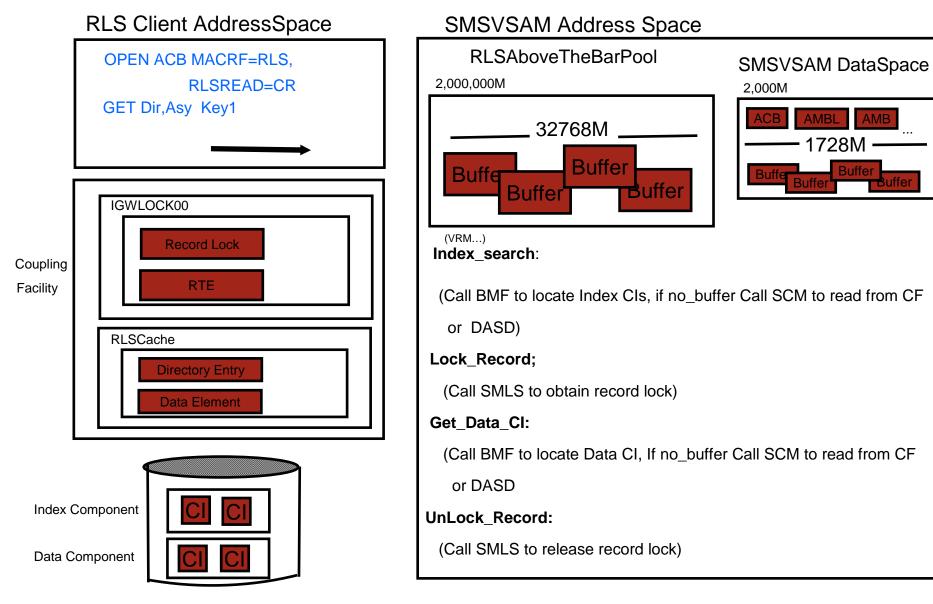


Control Interval



CICS3.Tran3

## **Overview of Get Path**



## **RLS - RAS**

- RLS provides extensive first time data capture for logic errors.
  - •Many "health checks" in the code which produce ABEND0F4 dumps to capture the problem at the earliest possible point.
  - •All mainline paths protected by recovery routines which force the data set to be closed in order to prevent damage to the data set.
    - •Initial recovery design terminated SMSVSAM.
    - •New recovery design marks data set as unusable.
  - •Extensive logging and tracing facilities.
  - •RAS is considered a high priority element of RLS design..
- End result:
  - Problems easier to debug..
  - •Much less likely for broken data sets or data integrity problems.

### **RLS Performance Measurements**

#### SMF 62 and 64

- SMF 62 Created by RLS OPEN for each ACB.
- SMF 64 Created by RLS EOV and CLOSE for each ACB. Stats are on an ACB level.
- SMF 42 Subtypes 15, 16, 17, 18, 19
  - Subytpe 15 RLS statistics by Storage Class
  - Subtype 16 RLS statistics by Data set
    - Must use V SMS,MONDS(spherename),ON to collect subtype 16 statistics.
  - **Subtype 17** RLS locking Statistics for IGWLOCK00
  - **Subtype 18** RLS caching Statistics
  - Subtype 19 BMF statistics
- SMF formatter soon to be available as part of our IPCS VERBX SMSXDATA
- Note: Only one system in the sysplex collects the SMF 42 records. The system collecting the records is displayed in the D SMS,SMSVSAM operator command.

# RLS/TVS Configuration Change

#### **Configuration Changes**

Update CFRM policy to define lock, cache, list, log structures.

•See DFSMSdfp Storage Administration Reference for sizing info.

Update SYS1.PARMLIB(IGDSMSxx) with RLS/TVS parameters.

•See MVS Initialization and Tuning.

Define new SHCDSs (Share Control Data Sets).

•See DFSMSdfp Storage Administration Reference.

Update SMS configuration for Cache Sets.

•See DFSMSdfp Storage Administration Reference.

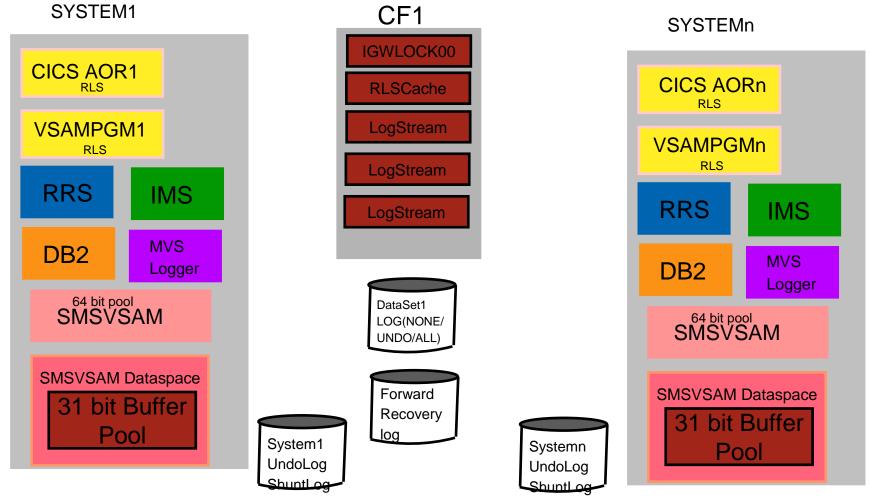
Update data sets with LOG(NONE/UNDO/ALL) and LOGSTREAMID.

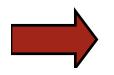
•See Access Methods Services for ICF.

## System Requirements - PARMLIB Changes

SYS1.PARMLIB(IGDSMSxx) SMS ACDS(acds) COMMDS(commds) INTERVAL(nnn|15) DINTERVAL(nnn|150) REVERIFY(YES|NO) ACSDEFAULTS(YES|NO) SYSTEMS(8|32) TRACE(OFF|ON) SIZE(nnnnnK|M) TYPE(ALL|ERROR) JOBNAME(jobname|\*) ASID(asid|\*) SELECT(event, event....) DESELECT(event,event....) DSNTYPE(LIBRARY|PDS) DSSTIMEOUT(nnn|0) RLSMAXCFFEATURELEVEL(A|Z) RLS MAX POOL SIZE(nnn|100) RLSINIT(NO|YES) SMF\_TIME(NO|YES) CF\_TIME(nnn|3600) BMF\_TIME(nnn|3600) CACHETIME(nnn|3600) DEADLOCK\_DETECTION(iii|15,kkk|4) RLSTMOUT(nnn|0) RLSAboveTheBarMaxPoolSIze(system,size) RLSFixedPoolSize(system.size) SYSNAME(sys1,sys2,...) TVSNAME(nnn1,nnn2....) MAXLOCKS(max|0,incr|0) TV\_START\_TYPE(WARM|COLD,WARM|COLD...) AKP(nnn|1000,nnn|1000) LOG\_OF\_LOGS(logstream) QTIMEOUT(nnn|300)

# SYSPLEX with SMSVSAM (and TVS) - Example





# **SMSVSAM** Initialization

## **SMSVSAM** Initialization

**IGW619I ACTIVE SHARE CONTROL DATA SET 209** SYS1.DFPSHCDS.ACTIVE2.VSPLXPK ADDED. **IGW619I SPARE SHARE CONTROL DATA SET 283** SYS1.DFPSHCDS.SPARE.VSPLXPK ADDED. IGW3211 Running Protocol 4 IXL014I IXLCONN REQUEST FOR STRUCTURE IGWLOCK00 313 WAS SUCCESSFUL. JOBNAME: SMSVSAM ASID: 0009 CONNECTOR NAME: SYSTEM1 CFNAME: FACIL01 IGW321I System Ordinal is 1 IGW453I SMSVSAM ADDRESS SPACE HAS SUCCESSFULLY 316 CONNECTED TO DFSMS LOCK STRUCTURE IGWLOCK00 IGW3211 No retained locks IGW321I 0 RLS Sphere Record Table Entries read IGW321I 0 RLS Sphere Record Table Entries deleted IGW3211 No Spheres in lost locks

### SMSVSAM Initialization (cont.)

IGW414I SMSVSAM SERVER ADDRESS SPACE IS NOW ACTIVE.

IGW467I DFSMS RLS\_MAX\_POOL\_SIZE PARMLIB VALUE SET DURING 354 SMSVSAM ADDRESS SPACE INITIALIZATION ON SYSTEM: SYSTEM1 CURRENT VALUE: 100 IGW467I DFSMS DEADLOCK\_DETECTION PARMLIB VALUE SET DURING 355 SMSVSAM ADDRESS SPACE INITIALIZATION ON SYSTEM: SYSTEM1 THIS SYSTEM IS OPERATING AS THE GLOBAL DEADLOCK PROCESSOR. CURRENT VALUE: 15 4

IGW467I DFSMS RLS\_MAXCFFEATURELEVEL PARMLIB VALUE SET DURING SMSVSAM ADDRESS SPACE INITIALIZATION ON SYSTEM: SYSTEM1 CURRENT VALUE: Z

# SMSMVSAM Initialization (with TVS) - (cont.)

SYSTEM1 05008 11:34:01.17 IGW467I DFSMS TVSNAME PARMLIB VALUE SET DURING 578

SMSVSAM ADDRESS SPACE INITIALIZATION ON SYSTEM:

SYSTEM1 TVSNAME: IGWTV001

SYSTEM1 05008 11:34:01.18 IGW467I DFSMS TRANSACTIONAL VSAM UNDO LOG PARMLIB VALUE SET

DURING SMSVSAM ADDRESS SPACE INITIALIZATION ON SYSTEM:

SYSTEM1 UNDO LOGSTREAM NAME:

IGWTV001.IGWLOG.SYSLOG

SYSTEM1 05008 11:34:01.18 IGW467I DFSMS TRANSACTIONAL VSAM SHUNT LOG PARMLIB VALUE SET

DURING SMSVSAM ADDRESS SPACE INITIALIZATION ON SYSTEM:

SYSTEM1 SHUNT LOGSTREAM NAME:

IGWTV001.IGWSHUNT.SHUNTLOG

#### System Requirements - SMSVSAM Initialization - Example SYSTEM1

SYSTEM1 05008 11:34:01.18 IGW467I DFSMS TRANSACTIONAL VSAM TVS\_START\_TYPE PARMLIB

VALUE SET DURING SMSVSAM ADDRESS SPACE INITIALIZATION

ON SYSTEM: SYSTEM1 TVSNAME VALUE: IGWTV001

CURRENT VALUE: WARM 1

SYSTEM1 05008 11:34:06.29 IGW860I TRANSACTIONAL VSAM HAS SUCCESSFULLY REGISTERED

WITH RLS

SYSTEM1 05008 11:35:36.63 IGW865I TRANSACTIONAL VSAM INITIALIZATION IS COMPLETE.

SYSTEM1 05008 11:35:36.65 IGW886I 0 RESTART TASKS WILL BE PROCESSED DURING

TRANSACTIONAL RESTART PROCESSING

SYSTEM1 05008 11:35:36.65 IGW866I TRANSACTIONAL VSAM RESTART PROCESSING IS COMPLETE.

# **SMSVSAM Commands**

## SMSVSAM Display Commands

D SMS[,	
[,CFCACHE(structurename *	) ]
[,CFLS	]
[,CFVOL(volid)	]
[,DSNAME(dsn){,WTOR}	]
[,JOB(jobname){,WTOR}	]
[,LOG({logstreamid ALL}{,W	FOR} ]
[,MONDS(specmask *)	]
[,SHCDS	]
[,SHUNTED,{SPHERE(sphe	re) UR({urid ALL}}{,WTOR}]
[,SMSVSAM[,ALL]	]

# SMSVSAM Display Commands (cont)

D SMS[,

[,TRANVSAM[,ALL][,ALLLOGS][,WTOR]

1

[,URID({urid|ALL}){,WTOR} ]

D SMS, SMSVSAM, DIAG (CONTENTION)

#### D SMS, SMSVSAM (example)

#### D SMS, SMSVSAM

DISPLAY SMS,SMSVSAM - SERVER STATUS SYSNAME: SYSTEM1 AVAILABLE ASID: 0033 STEP: SmsVsamInitComplete

DISPLAY SMS, SMSVSAM - JOB STATUS

SUBSYSTEMS CONNECTED: 1 BATCH: 1

DISPLAY SMS, SMSVSAM - LOCK TABLE STATUS (IGWLOCK00)

CONNECT STATUS:

SYSNAME: SYSTEM1 ACTIVE RSN: 02010407 RbldNotActive

COMPOSITE STATUS:

ORIGINAL STRUCTURE: NOT VOLATILEFAILURE ISOLATEDNEWSTRUCTURE: NOT VOLATILEFAILURE ISOLATED

STRUCTURE STATUS:

SYSNAME: SYSTEM1 Duplex

# System Requirements - SMSVSAM Displays

13.19.03 S	YSTEM1	d sm	s,tranvs	am				
13.19.04 S	YSTEM1	IEE9	321 023					
IGW800I 1:	3.19.04 DISP	LAY SM	IS,TRAN	ISACTIO	NAL VSAM			
DISPLAY S	SMS,TRANSA	ACTION	AL VSAN	/I - SER\	/ER STATUS	S		
System	TVSNAME	State	Rrs	#Urs	Start		AKP	QtimeOut
SYSTEM1	IGWTV001	ACTIVE	REG	0	WARM/W/	ARM	200	400
DISPLAY S	SMS,TRANSA	ACTION	AL VSAN	/ - LOGS	STREAM ST	ATUS		
LogStream	nName		S	tate	Туре	Conn	ect Stat	us
IGWTV00 <sup>2</sup>	1.IGWLOG.S	YSLOG	E	nabled	UnDoLog	Conr	nected	
IGWTV00 <sup>2</sup>	1.IGWSHUN	L.SHUN	TLOG E	nabled	ShuntLog	Conn	ected	

## SMSVSAM Vary Commands

V SMS,{CFCACHE(cachename),{ENABLE|E}} {QUIESCE|Q} } { {CFVOL(volid),{ENABLE|E}} {QUIESCE|Q} { } {MONDS(dsname[,dsname...]),{ON|OFF} } {SHCDS(shcdsname),{NEW } } {NEWSPARE} } { {DELETE } } { {SMSVSAM,{ACTIVE } {FALLBACK } { {TERMINATESERVER { } {FORCEDELETELOCKSTRUCTURE { }

## SMSVSAM Vary Commands

```
V SMS,{TRANVSAM({tvsname|ALL}){,{QUIESCE|Q}}
                                                  }
             {,{ENABLE|E }}
                              }
{
             {,{DISABLE|D}} }
{LOG(logstreamid){{,QUIESCE|Q}}
                                      }
         {,{ENABLE|E }}
                               }
{
         {,{DISABLE|D}}
                         }
{SMSVSAM,SPHERE(sphere){,{QUIESCE|Q}}
                                             }
             {,{ENABLE|E }}
                              }
{TRANVSAM(tvsname), PEERRECOVERY{, {ACTIVE|A }}}
                 {,ACTIVEFORCE }}
{
                 {,{INACTIVE|I}}
```

# **RLS/CICS Environment**

#### **RLS/CICS** Environment

CICS and base VSAM FOR configuration.

•Advantages and disadvantages of the FOR/AOR configuration.

CICS and RLS configuration.

•Advantages and disadvantages of the CICS/RLS configuration.

RLS/CICS data recovery.

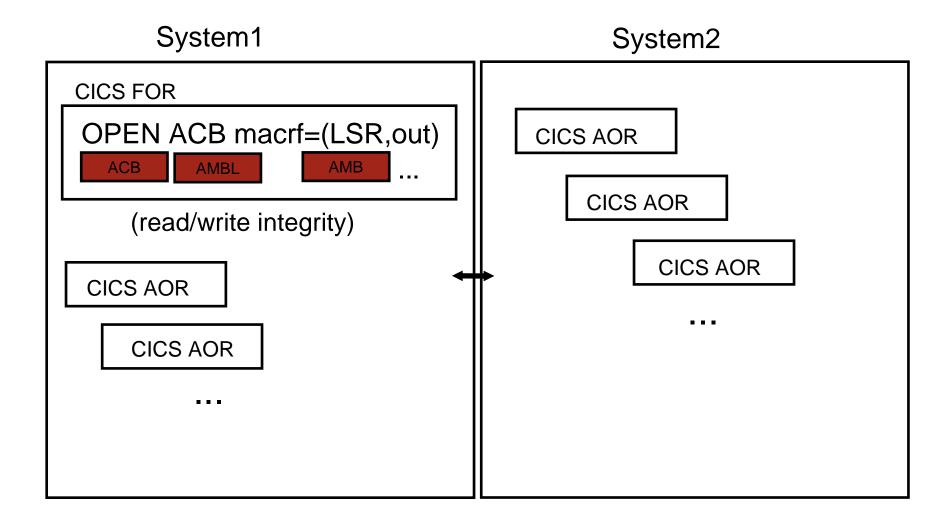
- Recoverable data sets.
- Recoverable subsystems.
- •Retained locks.
- •Lost locks.
- IDCAMS SHCDS commands
- QUICOPY/QUIBWO interface.

## **RLS/CICS** Environment

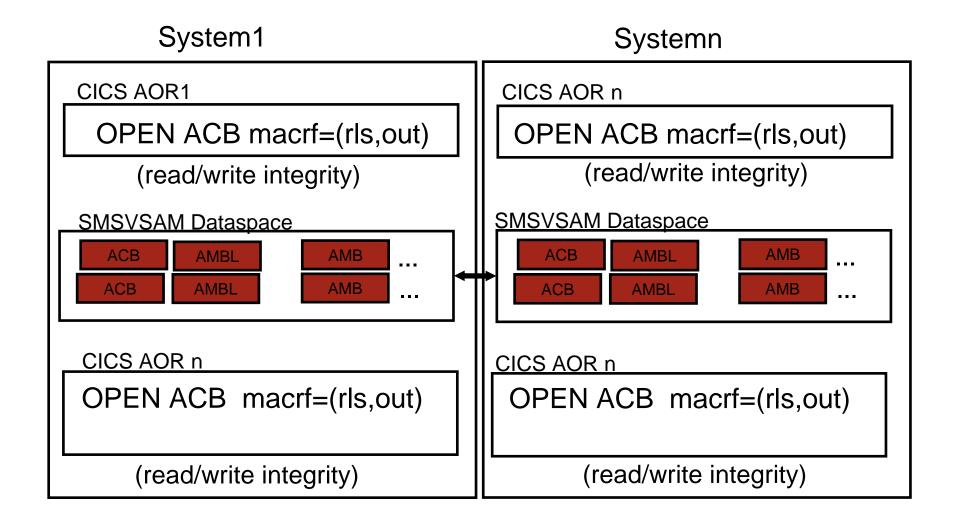
RLS/CICS automation enhancements.

•QUIOPEN/QUICLOSE interface.

## **CICS FOR/AOR Configuration**



## **RLS/CICS** Configuration



#### **RLS/CICS** Data Recovery

#### Recoverable data sets

- defined as LOG(UNDO/ALL) in the catalog.
  - UNDO backout logging performed by CICS (or TVS).
  - •ALL both backout and forward recovery logging (or TVS).
- LOG(ALL) data sets must have a LOGSTREAMID(forwardecoverylog) also defined in the catalog.

#### Non-Recoverable data sets

- defined as LOG(NONE) in the catalog.
  - No logging performed by CICS (or TVS).
- Recoverable Subsystems.
  - CICS (and TVS) must register with the SMSVSAM address space with a "subsystemname" so that locks obtained by that subsystem can be tracked.

#### **RLS/CICS** Data Recovery

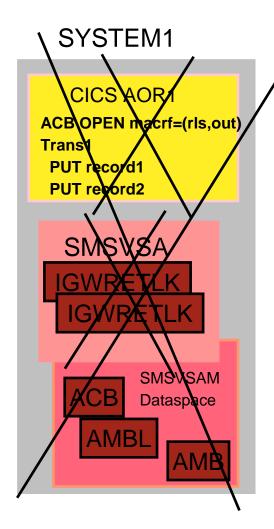
#### Retained locks

- Record locks are converted to "retained" in the event of a failure. The "owning" subsystem is the only subsystem that may access the record locks during recovery. All other subsystems or VSAM RLS applications will received a retained lock error in the RPL
- •SMSVSAM automatically notifies CICS when SMSVSAM restarts. CICS will automatically perform backouts when the file is reopened.

#### Lost Locks

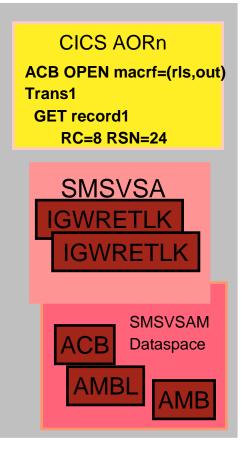
- •A data set which had actively held locks and a system failure occurs resulting in the loss of the RLS lock structure and at least one of the RLS address spaces at the exact same time.
- •Only the owning subsystem of the active locks may open the file and recovery the record locks. All other RLS opens will be failed until the data set has been fully recovered.

## **Retained Lock Example**



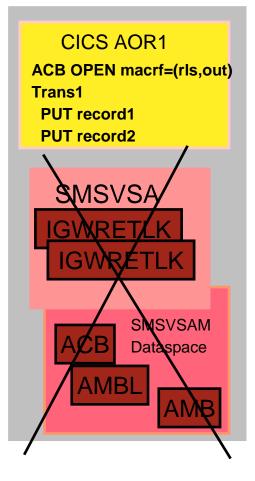
CF1								
IGWLOCK00								
Lock Table								
Record lock 1								
Record lock 2								
Record Table								
RTE lock 1 - (retained)								
RTE lock 1 - (retained)								
DataSet1 LOG(ALL) CICS logs								

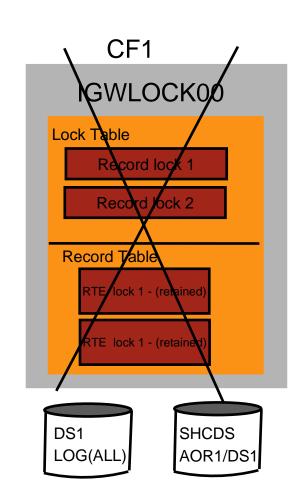
#### SYSTEMn



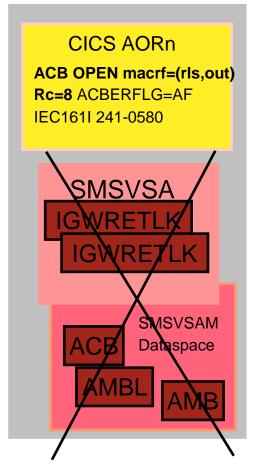
#### Lost Lock Example

#### SYSTEM1





#### SYSTEMn



#### **RLS/CICS** Data Recovery

#### IDCAMS SHCDS commands

•Used to list information about data set, clients, subsystems, etc. using RLS.

#### QUICOPY/QUIBWO interface.

- •Called by DSS to communicate with CICS (via the SMSVSAM) address space to inform CICS when a DSS copy/backup begins and ends.
- •Allows DSS to either take a "sharp" copy (via the QUICOPY interface) or a "fuzzy" copy (via the QUIBWO interface).
- •CICS will halt new transactions when a QUICOPY is under way. New opens will not be allowed during a QUICOPY.
- •CICS will log the start and end of the copy/backup operation. The data set can then be fully recovered from the last backup.

## SHCDS Commands

SHCDS {{LISTDS(base\_cluster\_name) {JOBS}} |

{LISTSUBSYS(subsystem\_name|ALL)} |

{LISTSUBSYSDS(subsystem\_name)} |

{LISTRECOVERY(base\_cluster\_name|ALL)} |

{LISTALL} |

{FRSETRR(base\_cluster\_name)} |

{FRUNBIND(base\_cluster\_name)} |

{FRBIND(base\_cluster\_name)} |

{FRRESETRR(base\_cluster\_name)} |

{FRDELETEUNBOUNDLOCKS(base\_cluster\_name)} |

{PERMITNONRLSUPDATE(base\_cluster\_name)} |

{DENYNONRLSUPDATE(base\_cluster\_name)} |

{REMOVESUBSYS(subsystem\_name)} |

{CFREPAIR({INFILE(ddname) |

INDATASET(datasetname)}

#### SHCDS Commands (continued)

{LIST|NOLIST})}

{CFRESET({INFILE(ddname) |

INDATASET(datasetname)}

{LIST|NOLIST})}

{CFREPAIRDS({base\_cluster\_name |

{partially\_qualified\_base\_cluster\_name)

{CFRESETDS({base\_cluster\_name |

{partially\_qualified\_base\_cluster\_name)

{LISTSHUNTED {SPHERE(base\_cluster\_name) |

URID(urid)

DATA(urid)}}

{RETRY {SPHERE(base\_cluster\_name) |

URID(urid)}}

{PURGE {SPHERE(base\_cluster\_name) |

URID(urid)}}

# SHCDS Example

**ISPF** Command Shell

Enter TSO or Workstation commands below:

===> SHCDS LISTS	SUBSYS(aor1)						
LISTING FROM SH	CDS IDCSH03						
		RECOVERY	LOCKS	LOCKS	LOCKS		
SUBSYSTEM NAME	STATUS	NEEDED	HELD	WAITING	RETAINED		
AOR1	ONLINEFAILED	YES	0	0	1		
DATA SETS IN LOST LOCKS 0							
DATA SETS IN NON-RLS UPDATE STATE 0							
TRANSACTION COL	TRANSACTION COUNT 1						
ala ala ala							

\*\*\*

# SHCDS Example

**ISPF** Command Shell

Enter TSO or Workstation commands below:

===> SHCDS LISTDS('dataset1\*')

----- LISTING FROM SHCDS ----- IDCSH02 ------

DATA SET NAME----dataset1

CACHE STRUCTURE----CACHE01

RETAINED LOCKS------YES NON-RLS UPDATE PERMITTED-----NO

LOST LOCKS------NO PERMIT FIRST TIME-----NO

LOCKS NOT BOUND------NO FORWARD RECOVERY REQUIRED-----NO

RECOVERABLE-----YES

# SHCDS Example (cont.)

#### SHARING SUBSYSTEM STATUS SUBSYSTEM SUBSYSTEM RETAINED LOST NON-RLS UPDATE NAME STATUS LOCKS LOCKS PERMITTED ---------------------------AOR1 ONLINE--FAILED YES NO NO

\*\*\*

# RLS/CICS Automation Enhancements

### QUIOPEN/QUICLOSE Interface

•QUICLOSE interface is used by CICS to fully close a data set around the sysplex.

•SMSVSAM drives CICS quiesce exit which issues closes for all regions open to the data set.

•SMSVSAM updates the catalog and marks the data set as quiesced.

•RLS opens against a quiesced data set will be failed.

•QUIOPEN interface is used by CICS to enable a data set to be reopened for RLS use.

•SMSVSAM drives CICS quiesce exit to ALL CICS regions registered with RLS.

•SMSVSAM updates the catalog and marks the data set as unquiesced.

Invoked with the following commands:

•V SMS,SMSVSAM,SPHERE(spherename),Q

•V SMS,SMSVSAM,SPHERE(spherename),E

•F cicsname,CEMT SET DSN(RLSADSW.VFA1D.\*),QUI

•F cicsname,CEMT SET DSN(RLSADSW.VFA1D.\*),UNQ

# **Transactional VSAM (TVS)**

# Transactional VSAM (TVS)

- Enhance VSAM Record Level Sharing (RLS) to provide data recovery capabilities for any application exploiting VSAM RLS.
- VSAM RLS data recovery capabilities include:
  - transactional recovery
  - data set recovery
- VSAM RLS becomes a "transactionalized" access method, or is now referred to as "Transactional VSAM" (TVS).

# System Requirements -Hardware/Software Requirements

- Parallel sysplex running z/OS 1.4 or higher with VSAM RLS implemented.
- z/OS Transactional VSAM (separately priced feature).
- z/OS RRMS implemented.
- z/OS System Logger implemented.
- CICS VSAM Recovery (CICVR) Utility (optional)

# Application Requirements - Data Set Changes

- Data sets accessed by RLS must have a LOG parm specifed in the catalog. Valid values are:
  - •LOG(NONE) Non-recoverable data set. Can be opened for input/output by any RLS application.
  - •LOG(UNDO) Recoverable data set requiring backout (UNDO) logging. Can be opened for input/output by RLS recoverable subsystems (i.e. CICS) and/or RLS applications running on a z/OS system with the TVS feature installed.
  - •LOG(ALL) Recoverable data set requiring both backout (undo) and forward recovery logging. Can be opened for input/output by RLS recoverable subsystems (i.e. CICS) and/or RLS applications running on a z/OS system with the TVS feature installed.

# Application Requirements - Data Set Changes (cont)

 Data sets defined as LOG(ALL) must also have a LOGSTREAMID(fowardrecoverylogname) specified in the catalog.

# Application Requirements - Data Set Define/Alter Example

DEFINE CLUSTER (NAME(recoverabledataset)

RECORDSIZE(100 100)

STORCLAS(storclasname)

FSPC(20 20)

LOG (ALL) -

SHAREOPTIONS(2 3)

LOGSTREAMID(forwardrecoverylog)

-

CISZ(512)

KEYS(06 8) INDEXED

)

DATA(NAME(recoverabledataset.DATA) -

VOLUME(volser)

TRACKS (1,1)) -

INDEX(NAME(recoverabledataset.INDEX) -

VOLUME(volser)

TRACKS (1,1))

# Application Requirements – RLS/TVS Access Options

- Transactional VSAM support occurs when:
  - ACB MACRF=(RLS,OUT) for recoverable data set (LOG(UNDO|ALL))
  - ◆ACB MACRF=(RLS,IN), RLSREAD=CRE .
  - //ddname DD DSN=recoverabledatasetname,DISP=shr,RLS=(CR|NRI) and ACB MACRF=(OUT)
  - //ddname DD DSN=datasetname,DISP=shr,RLS=CRE and ACB MACRF=(IN)

# Application Requirements -Transactional Recovery

- RLS applications opening recoverable data sets on z/OS with the TVS feature installed, <u>should</u> be modified to add SRRCMIT and SRRBACK interfaces.
- SRRCMIT and SRRBACK will either commit or backout the unit of recovery (UR) provided by SMSVSAM on behalf of the VSAM RLS application.
- Explicitly committing or backing out the UR will release record level locks in a timely fashion. Failure to do so may impact other sharers of the data set.
- SMSVSAM will implicitly issue a commit or backout at EOT, if the VSAM application fails to do so.

# Application Requirements -Supported Languages

High level language support for RLS and RRS interfaces:

•PLI

•C & C++

•COBOL

Assembler

# Application Requirements - Explicit Commit Example

//ddname DD DSN=Recoverabledatasetname,DISP=SHR				
//step1 EXEC PGM=vsamrlspgm				
Begin JOB Step No lo	ocks held			
OPEN ACB MACRF=(RLS,OUT)				
(UR1)				
GET UPD record 1 Obta	ain an exclusive lock on record 1			
PUT UPD record 1 Loci	k on record 1 remains held			
GET repeatable read record n Obta	in a shared lock on record n			
PUT ADD record n+1 Obta	in an exclusive lock on record n+1			
GET UPD record 2 Obtain	in an exclusive lock on record 2			
PUT UPD record 2 Lock	on record 2 remains held			
Call SRRCMIT Con	nmit changes, all locks released .			
CLOSE				

End of JOB Step

# Application Requirements - Implicit Commit Example

//ddname DD DSN=Recoverabledatasetname,DISP=SHR				
//step1 EXEC PGM=vsamrlspgm				
Begin JOB Step No locks held				
OPEN ACB MACRF=(RLS,OUT)				
(UR1)				
GET UPD record 1 Obtain an exclusive lock on record 1				
PUT UPD record 1 Lock on record 1 remains held				
GET repeatable read record n Obtain a shared lock on record n				
PUT ADD record n+1 Obtain an exclusive lock on record n+1				
GET UPD record 2 Obtain an exclusive lock on record 2				
PUT UPD record 2 Lock on record 2 remains held				
CLOSE All Locks are retained				
End of JOB Step (normal) Commit changes release all locks				

# Application Requirements - Explicit Backout Example

//ddname DD DSN=Recoverabledatasetname,DISP=SHR				
//step1 EXEC PGM=vsamrlspgm				
Begin JOB Step	No locks held			
OPEN ACB MACRF=(RLS,OUT)				
(UR1)				
GET UPD record 1	Obtain an exclusive lock on record 1			
PUT UPD record 1	Lock on record 1 remains held			
GET repeatable read record n	Obtain a shared lock on record n			
PUT ADD record n+1	Obtain an exclusive lock on record n+1			
GET UPD record 2	Obtain an exclusive lock on record 2			
PUT UPD record 2	Lock on record 2 remains held			
Call SRRBACK	Undo changes, all locks released .			
CLOSE				

End of JOB Step

# Application Requirements - Implicit Backout Example

//ddname DD DSN=Recoverabledatasetname,DISP=SHR				
//step1 EXEC PGM=vsamrlspgm				
Begin JOB Step	No locks held			
OPEN ACB MACRF=(RLS,OUT)				
(UR1)				
GET UPD record 1	Obtain an exclusive lock on record 1			
PUT UPD record 1	Lock on record 1 remains held			
GET repeatable read record n	Obtain a shared lock on record n			
PUT ADD record n+1	Obtain an exclusive lock on record n+1			
GET UPD record 2	Obtain an exclusive lock on record 2			
PUT UPD record 2	Lock on record 2 remains held			
Cancel				

End of JOB Step (abnormal) ------ Undo changes release all locks

# **Information about TVS**

#### Information about DFSMS and TVS

•www.storage.ibm.com/software/sms/index.html

•www.storage.ibm.com/software/sms/tvs/index.html

**Additional Information** 

#### •www.redbooks.ibm.com

•	Transactional VSAM Presentation Guide	SG24-6973
•	Transactional VSAM Overview and Planning Guide	SG24-6971
•	Transactional VSAM Application Migration Guide	SG24-6972
•	VSAM Demystified	SG24-6105

# Multiple Lock Structure (MLS)

# **Multiple Lock Structure**

- Multiple Lock Structures (MLS), goal of this function is to remove the single point of failure of one lock structure in the current VSAM RLS design
  - Current Locking Design
  - Current Locking Design Issues
  - Multiple lock Structure Design

# **Current Locking Design**

- The current design of locking uses one coupling facility (CF) lock structure,
   IGWLOCK00, which contains:
  - Record locks and record data (retained locks)
  - System "Special" locks:
    - Sphere, component, subsystem locks and data set related record data

# **Current Locking Design - Issues**

- □ The current locking design has two issues:
  - IGWLOCK00 represents a single point of failure in the sysplex:
    - A "run away" application could fill IGWLOCK00 with record locks, causing all RLS application's lock requests in the sysplex to fail.
  - IGWLOCK00 could cause performance issues:
    - All RLS locking activity against a single lock structure in a single CF

# **Proposed Design - Multiple Lock Structure**

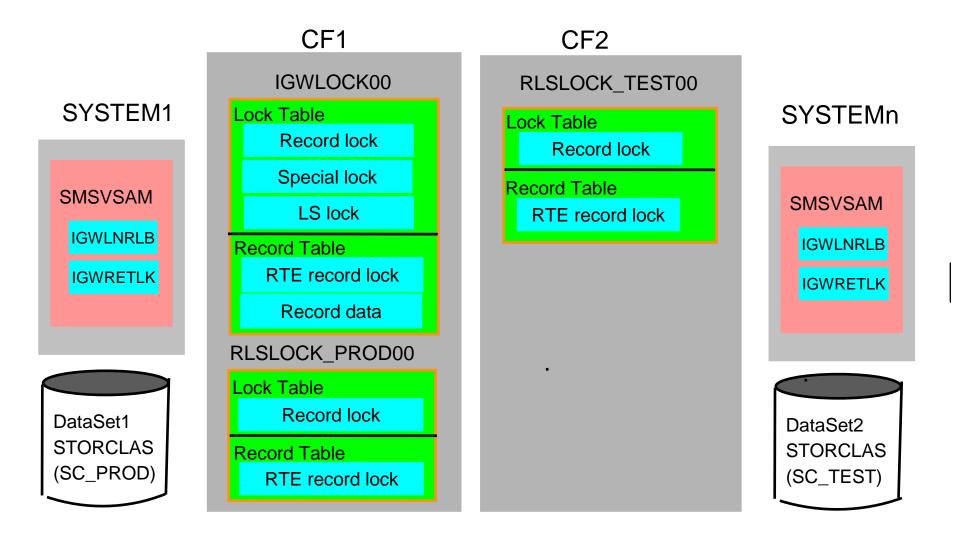
# □ Continue to support IGWLOCK00 as the "primary" lock structure, which will contain:

- Record locks and record data for data sets not using the new MLS support
- System "special" locks:
  - Sphere, component, subsystem locks and data set record data.
- "Lock structure" lock (associates data sets to lock structures)

#### □ Add new "secondary" lock structures, which will contain:

- Record locks and record data for data sets using the new MLS support
- □ Assign data sets to "secondary" lock structures via a new "lock set" parameter on the SMS STORCLAS construct
  - A "secondary" lock structure will be assigned from the list of lock structures specified in the lock set parameter
  - If the lock set parameter is blank, IGWLOCK00 will be assigned as the default

# Mulitple Lock Structure Example



•OA21101

•D SMS,SMSVSAM,QUIESCE

OA19421

Move index buffer above the bar for release 1.7 and above

OA19975

Change the wait time for the castout lock in the RLS read path from 0.03second/0.000026 seconds to 0.0015 seconds.

OA16676, OA16870, OA17643

Remove Assignedspheres ENQ hang

OA17644, OA18070, OA18541, OA18285, OA18688,
 OA18902

•SCM RAS APARs

OA20367

RLS/Catalog hang in Open/Delete

OA21705

•Fix the storage leaks in MMFSTUFF dataspace

OA18933

SSF compress/expand pool failure

### OA17556

•D SMS,SMSVSAM,DIAG(CONTENTION)

- Display TCBs in latches contention
- OA12045, OA12851, OA16982
  - VERBX IGWFPMAN 'F(IPCS)' From IPCS Panel
    - •Q Analyze current Failure
    - AS Analyze current Address Space Threads
    - POOLS Analyze SSF Pools

D SMS, SMSVSAM, DIAG(CONTENTION) - example #1

SYSTEM1	d sms	,smsvsam,o	diag(cor	ntention)				
SYSTEM1	IGW3	IGW343I VSAM RLS DIAG STATUS (V.01)						
RESOUR	CE		WAI	TER	HOLD	ER	ELAPSED	
TYPE	ID JOI	B NAME	ASID	TASK	ASID	TASK	TIME	
LATCH 7F	158C70 SI	MSVSAM	003A	008DA2	50 003	A 008D	7218 00:00:0	6
DESCRIPTION <sup>,</sup> IGWI YSPH - SHM OBJECT POOL								

DESC	RIPTION. IGVIL13PT - 3F		ECTPUUL		
LATCH	7F151E78 SMSVSAM	003A	008D7218	003A 008DC1C8 00:00:21	
DESCRIPTION: IGWLYDTS - SHM OBJECT POOL					
LATCH	7BAD43B8 SMSVSAM	003A	008DC1C8	002D 007F3000 00:19:09	
LATCH	7BAD43B8 SMSVSAM	003A	008D5A48	002D 007F3000 00:22:09	
LATCH	7BAD43B8 SMSVSAM	003A	008D6938	002D 007F3000 00:33:23	
LATCH	07F1B1D0 SMSVSAM	003A	008D64F8	003A 008D6CF0 01:47:20	
LATCH	07F1D3B8 SMSVSAM	003A	008D6CF0	0000 0000000 11:23:30	

D SMS, SMSVSAM, DIAG (CONTENTION) - example #2

SYSTEM1	d sms,smsvsam,	,diag(contention)			
SYSTEM1	IGW343I VSAM RLS DIAG STATUS (V.01)				
RESOUR	CE	WAITER   HOLDER  ELAPSED			
TYPE	ID JOB NAME	ASID TASK ASID TASK TIME			
LATCH 7E	BAD43B8 SMSVSAM	003A 008D5A48 003A 007F3000 00:22	:09		
LATCH 07	F1B1D0 SMSVSAM	003A 007F3000 003A 008D5A48 00:22	2:09		
LATCH 07	F1B1D0 SMSVSAM	003A 008D64F8 003A 008D5A48 00:22	:24		
LATCH 07	F1B1D0 SMSVSAM	003A 008D6CF0 003A 008D5A48 00:23	3:30		

IP VERBX IGWFPMAN 'F(IPCS)' - example

Function(F) Component AddressSpace Analysis IPCSPrint Help -----SMS PDSE IPCS MAIN------

COMMAND===>

Function( ) Component( ) CB@(0000000 )

JOB(SMSVSAM) or ASID(000A)

VERB===> IGWFPMAN

Primary(000A : SMSVSAM ) Secondary(000A : SMSVSAM )

Dump: Dump Name Title: Dump Title

# Summary

- RLS provides full read/write integrity to your existing VSAM files.
- RLS can improve both performance and availability in your CICS and non-CICS VSAM environments.
- RLS provides data protection after a system failure.
- RLS provides automation for data recovery.
- Improved RAS
- Minimal application/configuration changes required.

# Summary

RLS has been enhanced to perform data recovery in the form of:

- transactional recovery
- data set recovery.
- VSAM RLS Applications can take advantage of RLS's new data recovery by using the RRS commit and backout protocols.
- VSAM RLS Applications should reconsider restart procedures in a shared environment.

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